

# Executable Semantics for D-Clean

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# Overview

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# D-Clean at a glance

- A Clean-like language with distributed computations over a cluster
- Preserves the functional style
- Tasks are pure functional computations
- The dataflow is controlled by the coordination language elements
- Common algorithmic skeletons can be written elegantly using the newly defined language elements
- A D-Clean program consists of a start expression = composition of coordination structures parameterized by Clean expressions

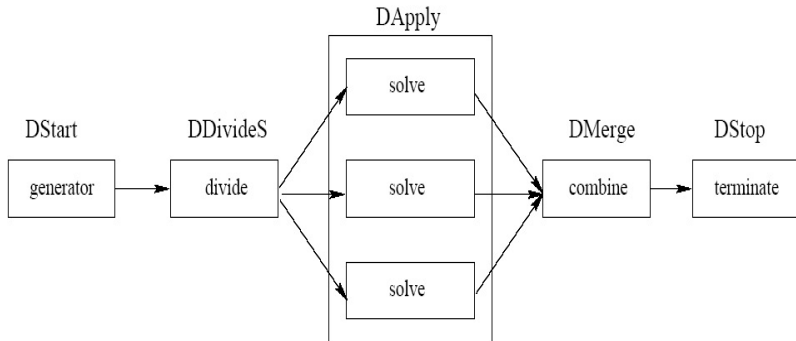


# Execution model

- The distributed computation is described using D-Clean primitives
- The primitives are compiled into D-Box expressions
- D-Box primitives are generating Clean computation nodes
- The nodes are communicating via channels
- The middleware will distribute the nodes over the cluster, where they will be started as individual Clean programs



# Small example

$$\text{DistrStart} = (\text{DStop terminate}) \circ (\text{DMerge mergesort}) \circ$$
$$(\text{DApply quicksort}) \circ (\text{DDivideS divide 3}) \circ$$
$$(\text{DStart [1..100000]})$$


# Why executable version?

- the well-formedness of the semantics can be checked by the Clean compiler
- the semantics can be executed
- properties of the semantics can be tested automatically
- visualization can be generated showing the parallelism automatically



## BNF

Instead of an ugly BNF description:

```
⟨DISTART_RULE⟩ ::= "DistrStart" "=" ⟨DEXPR⟩
⟨DEXPR⟩         ::= ⟨DPRIMITIVE⟩ | ⟨SCHEME_NAME⟩ { ⟨act_param⟩ }*
                  | ⟨DEXPR⟩ ⟨DEXPR⟩
⟨DPRIMITIVE⟩    ::= ⟨DStart_USE⟩ | ⟨DStop_USE⟩ | ⟨DMap_USE⟩ |
                  | ⟨DDivideS_USE⟩ | ⟨DMerge_USE⟩
⟨SCHEME_DEF⟩    ::= "SCHEME" ⟨SCHEME_NAME⟩ { ⟨formal_param⟩ }*
                  "==" { ⟨DEXPR⟩ }+
⟨SCHEME_NAME⟩  ::= { ⟨UpCaseLetter⟩ }+
```

we provide Clean datatypes and executable programs.



# Used types for executable semantics

*// either a function or a composition of functions*

```
:: DExpr a b = F String (a → b)
             | D (DFun (Ch a) (Ch b))
```

*// a DClean expression is a state transition function*

```
:: DFun a b ::= a State → (b, State)
```

*// single channel*

```
:: Ch a
```

*// multiple channel*

```
:: MCh a ::= [Ch a]
```

*// wrapper*

```
DExec :: a (DFun a b) → (a, (b, State))
```

```
DExec a expr = (a, DStop (DStart a expr (0, [])))
```





# D-Clean primitives

`DApply` :: `String (a → b) → DFun (Ch a) (Ch b) | toString b`

`DDivideS` :: `(Int [a] → [[a]]) Int → DFun (Ch [a]) (MCh [a])`  
| `toString a`

`DDivideD` :: `(a → [b]) → DFun (Ch a) (MCh b) | toString b`

`DMerge` :: `([b] → a) → DFun (MCh b) (Ch a) | toString a`



## Farm

Start `w = dumpGraph (DExec myval myflow) w`

where

`myval = [1..10]`

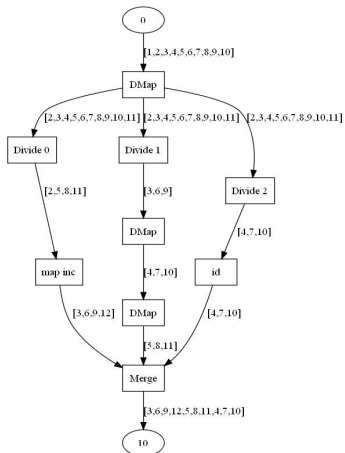
`myflow i`

`= DMap inc (0,i)`

`>>= DDivideD (divide 3)`

`>>= DApplyN [F "map inc" (map inc),  
D ( $\lambda i \rightarrow$  DMap inc i  $>>=$  DMap inc),  
F "id" id]`

`>>= DMerge flatten`



# D-Clean primitives

```
DApplyN :: [DExpr a b] → DFun (MCh a) (MCh b) | toString b
```

```
DApplyN funs = handleDExpr funs
```

**where**

```
handleDExpr :: [DExpr a b] (MCh a) State → (MCh b, State)
              | toString b
```

```
handleDExpr [] [] state = ([], state)
```

```
handleDExpr [] inflows state = abort "run-time error"
```

```
handleDExpr [F name fun:funs] [i:is] state
```

```
# (result,state) = mkFunBox name fun i state
```

```
# (results,state) = handleDExpr funs is state
```

```
= ([result:results],state)
```

```
handleDExpr [D dfun:funs] [i:is] state
```

```
# (result,state) = dfun i state
```

```
# (results,state) = handleDExpr funs is state
```

```
= ([result:results],state)
```



# Utility primitives

$DMap :: (a \rightarrow b) \rightarrow DFun (Ch [a]) (Ch [b]) \mid toString\ b$

$DReduce :: ([b] \rightarrow a) \rightarrow DFun (MCh [b]) (MCh a) \mid toString\ a$

$DProduce :: (a \rightarrow [b]) \rightarrow DFun (MCh a) (MCh [b]) \mid toString\ b$

$DFilter :: (a \rightarrow Bool) \rightarrow DFun (MCh [a]) (MCh [a]) \mid toString\ a$



# Code transformations

- optimization on the amount of communications on the channels
- change on the granularity of the computations

$$\begin{aligned} & \text{DApplyN } [F \text{ "f1" } f1, \dots, F \text{ "fn" } fn] \\ & \gg= \text{DApplyN } [F \text{ "g1" } g1, \dots, F \text{ "gn" } gn] \\ \equiv & \text{DApplyN } [F \text{ "fg1" } (g1 \circ f1), \dots, F \text{ "fgn" } (gn \circ fn)] \end{aligned}$$



# Properties

For all lists and appropriate functions f:

```
map f list == DDivideD (divide n) (0 ,list)
             >>= DApply1 (F "f" f)
             >>= DMerge id
```

To test the correctness of the transformations: for all channels c and appropriate functions f and g:

```
(DApply1 (F "f" f) >>= DApply1 (F "g" g)) c
== DApply1 (F "f" (g o f)) c
```



## Conclusions and future work

- Executable semantics for the D-Clean primitives
- To mimic the functionality of the distributed programs
- Sanity check of the semantics
- Transformations and properties - decreasing the workload of the channels
- Visualizing parallelism
- Examples

